# python-nnf

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python-nnf is a package for working with logical sentences written in the negation normal form.

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	Installing
python-nnf can be installed with pip:	
pip installuser nnf	

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## CHAPTER 2

### Module reference

### 2.1 Module contents

```
class nnf.NNF
```

Bases: object

Base class for all NNF sentences.

```
clause() \rightarrow bool
```

The sentence is a clause.

Clauses are Or nodes with variable children that don't share names.

```
condition (model: Dict[Hashable, bool]) → nnf.NNF
```

Fill in all the values in the dictionary.

```
consistent() \rightarrow bool
```

Some set of values exists that makes the sentence correct.

This method doesn't necessarily try to find an example, which can make it faster. It's decent at decomposable sentences and sentences in CNF, and bad at other sentences.

```
contradicts (other: nnf.NNF) \rightarrow bool
```

There is no set of values that satisfies both sentences.

```
decision\_node() \rightarrow bool
```

The sentence is a valid binary decision diagram (BDD).

```
\texttt{decomposable}\,(\,)\,\to bool
```

The children of each And node don't share variables, recursively.

```
\texttt{deduplicate()} \rightarrow T\_NNF
```

Return a copy of the sentence without any duplicate objects.

If a node has multiple parents, it's possible for it to be represented by two separate objects. This method gets rid of that duplication.

It's better to avoid the duplication in the first place. This method is for diagnostic purposes, in combination with object\_count().

#### $deterministic() \rightarrow bool$

The children of each Or node contradict each other.

May be very expensive.

#### entails (other: nnf.NNF) $\rightarrow$ bool

Return whether other is always true if the sentence is true.

This is faster if self is a term or other is a clause.

#### equivalent (other: nnf.NNF) $\rightarrow$ bool

Test whether two sentences have the same models.

If the sentences don't contain the same variables they are considered equivalent if the variables that aren't shared are independent, i.e. their value doesn't affect the value of the sentence.

```
flat() \rightarrow bool
```

A sentence is flat if its height is at most 2.

That is, there are at most two layers below the root node.

#### $forget(names: Iterable[Hashable]) \rightarrow nnf.NNF$

Forget a set of variables from the theory.

Has the effect of returning a theory without the variables provided, such that every model of the new theory has an extension (i.e., an assignment) to the forgotten variables that is a model of the original theory.

**Parameters** names – An iterable of the variable names to be forgotten

#### $forget_aux() \rightarrow nnf.NNF$

Returns a theory that forgets all of the auxillary variables

```
\textbf{height} () \rightarrow int
```

The number of edges between here and the furthest leaf.

#### $\textbf{implicants} () \rightarrow nnf.Or[nnf.And[nnf.Var]][nnf.And[nnf.Var]]$

Extract the prime implicants of the sentence.

Prime implicants are the minimal terms that imply the sentence. This method returns a disjunction of terms that's equivalent to the original sentence, and minimal, meaning that there are no terms that imply the sentence that are strict subsets of any of the terms in this representation, so no terms could be made smaller.

#### $implicates() \rightarrow nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]]$

Extract a prime implicate cover of the sentence.

Prime implicates are the minimal implied clauses. This method returns a conjunction of clauses that's equivalent to the original sentence, and minimal, meaning that there are no clauses implied by the sentence that are strict subsets of any of the clauses in this representation, so no clauses could be made smaller.

While implicants () returns all implicants, this method may only return some of the implicates.

#### $implies(other: nnf.NNF) \rightarrow bool$

Return whether other is always true if the sentence is true.

This is faster if self is a term or other is a clause.

#### $is\_CNF (strict: bool = False) \rightarrow bool$

Return whether the sentence is in the Conjunctive Normal Form.

**Parameters strict** – If True, follow the definition of the Knowledge Compilation Map, requiring that a variable doesn't appear multiple times in a single clause.

#### is DNF (strict: bool = False) $\rightarrow bool$

Return whether the sentence is in the Disjunctive Normal Form.

**Parameters strict** – If True, follow the definition of the Knowledge Compilation Map, requiring that a variable doesn't appear multiple times in a single term.

#### is MODS() $\rightarrow$ bool

Return whether the sentence is in MODS form.

MODS sentences are disjunctions of terms representing models, making the models trivial to enumerate.

#### $leaf() \rightarrow bool$

True if the node doesn't have children.

That is, if the node is a variable, or one of true and false.

#### $make\_pairwise() \rightarrow nnf.NNF$

Alter the sentence so that all internal nodes have two children.

This can be easier to handle in some cases.

#### make smooth() $\rightarrow$ nnf.NNF

Transform the sentence into an equivalent smooth sentence.

#### ${\tt mark\_deterministic}\,(\,)\,\to None$

Declare for optimization that this sentence is deterministic.

Note that this goes by object identity, not equality. This may matter in obscure cases where you instantiate the same sentence multiple times.

#### marked deterministic() $\rightarrow$ bool

Whether this sentence has been marked as deterministic.

#### $model\_count() \rightarrow int$

Return the number of models the sentence has.

This can be done efficiently for sentences that are decomposable and deterministic.

#### $models() \rightarrow Iterator[Dict[Hashable, bool]]$

Yield all dictionaries of values that make the sentence correct.

Much faster on sentences that are decomposable. Even faster if they're also deterministic.

#### $negate() \rightarrow nnf.NNF$

Return a new sentence that's true iff the original is false.

#### $object\_count() \rightarrow int$

Return the number of distinct node objects in the sentence.

#### project (names: Iterable[Hashable]) → nnf.NNF

Dual of forget (): will forget all variables not given

#### $satisfiable() \rightarrow bool$

Some set of values exists that makes the sentence correct.

This method doesn't necessarily try to find an example, which can make it faster. It's decent at decomposable sentences and sentences in CNF, and bad at other sentences.

#### **satisfied\_by** (*model: Dict[Hashable, bool]*) → bool

The given dictionary of values makes the sentence correct.

#### $simplify (merge\_nodes: bool = True) \rightarrow nnf.NNF$

Apply the following transformations to make the sentence simpler:

• If an And node has false as a child, replace it by false

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- If an Or node has true as a child, replace it by true
- Remove children of And nodes that are true
- Remove children of Or nodes that are false
- If an And or Or node only has one child, replace it by that child
- If an And or Or node has a child of the same type, merge them

**Parameters merge\_nodes** – if False, don't merge internal nodes. In certain cases, merging them may increase the size of the sentence.

#### $simply_conjunct() \rightarrow bool$

The children of And nodes are variables that don't share names.

#### $simply_disjunct() \rightarrow bool$

The children of Or nodes are variables that don't share names.

```
size() \rightarrow int
```

The number of edges in the sentence.

Note that sentences are rooted DAGs, not trees. If a node has multiple parents its edges will still be counted just once.

#### $smooth() \rightarrow bool$

The children of each Or node all use the same variables.

#### $solve() \rightarrow Optional[Dict[Hashable, bool]]$

Return a satisfying model, or None if unsatisfiable.

```
term() \rightarrow bool
```

The sentence is a term.

Terms are And nodes with variable children that don't share names.

```
\textbf{to\_CNF} \ (simplify: bool = True) \ \rightarrow nnf. And [nnf. Or [nnf. Var]] [nnf. Or [nnf. Var]] [nnf. Var] [nnf.
```

Compile theory to a semantically equivalent CNF formula.

**Parameters** simplify – If True, simplify clauses even if that means eliminating variables.

```
to_DOT (*, color: bool = False, color_dict: Optional[Dict[str, str]] = None, label: str = 'text', label_dict: Optional[Dict[str, str]] = None) \rightarrow str
Return a representation of the sentence in the DOT language.
```

DOT is a graph visualization language.

#### **Parameters**

- **color** If True, color the nodes. This is a bit of an eyesore, but might make them easier to understand.
- label-If 'text', the default, label nodes with "AND", "OR", etcetera. If 'symbol', label them with unicode symbols like "" and "".
- color\_dict Use an alternative palette. This should be a dictionary with keys 'and', 'or', 'true', 'false', 'var' and 'neg'. Not all keys have to be included. Passing a dictionary implies color=True.
- label\_dict Use alternative labels for nodes. This should be a dictionary with keys 'and', 'or', 'true' and 'false'. Not all keys have to be included.

```
\textbf{to\_MODS} \ () \ \rightarrow nnf.Or[nnf.And[nnf.Var]][nnf.And[nnf.Var]][nnf.And[nnf.Var]]
```

Convert the sentence to a MODS sentence.

(continues on next page)

```
to model () → Dict[Hashable, bool]
           If the sentence directly represents a model, convert it to that.
           A sentence directly represents a model if it's a conjunction of (unique) variables, or a single variable.
      valid() \rightarrow bool
           Check whether the sentence is valid (i.e. always true).
           This can be done efficiently for sentences that are decomposable and deterministic.
      vars() → FrozenSet[Hashable]
           The names of all variables that appear in the sentence.
      walk() \rightarrow Iterator[nnf.NNF]
           Yield all nodes in the sentence, depth-first.
           Nodes with multiple parents are yielded only once.
class nnf.Internal(children: Iterable[T_NNF_co] = ())
      Bases: nnf.NNF, typing.Generic
      Base class for internal nodes, i.e. And and Or nodes.
      children
      leaf() \rightarrow bool
           True if the node doesn't have children.
           That is, if the node is a variable, or one of true and false.
      map(func: Callable[[T\_NNF\_co], U\_NNF]) \rightarrow nnf.Internal[\sim U\_NNF][U\_NNF]
           Apply func to all of the node's children.
class nnf.And(children: Iterable[T_NNF_co] = ())
      Bases: nnf.Internal
      Conjunction nodes, which are only true if all of their children are.
      decision node() \rightarrow bool
           The sentence is a valid binary decision diagram (BDD).
class nnf.Or(children: Iterable[T_NNF_co] = ())
      Bases: nnf.Internal
      Disjunction nodes, which are true if any of their children are.
      {\tt decision\_node}\,(\,)\,\to bool
           The sentence is a valid binary decision diagram (BDD).
class nnf.Var (name: Hashable, true: bool = True)
      Bases: nnf. NNF
      A variable, or its negation.
      If its name is a string, its repr will use that name directly. Otherwise it will use more ordinary constructor syntax.
      >>> a = Var('a')
      >>> a
      а
      >>> ~a
      ~a
      >>> b = Var('b')
      >>> a | ~b == Or({Var('a', True), Var('b', False)})
```

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True

>>> Var(10)

(continued from previous page)

```
Var(10)
     >>> Var(('a', 'b'), False)
      ~Var(('a', 'b'))
     static aux() \rightarrow nnf.Var
          Create an auxiliary variable with a unique label.
     decision node() \rightarrow bool
          The sentence is a valid binary decision diagram (BDD).
     name
     true
class nnf.Aux (hex=None, bytes=None, bytes_le=None, fields=None, int=None, version=None, *,
                   is_safe=<SafeUUID.unknown: None>)
     Bases: uuid.UUID
     Unique UUID labels for auxiliary variables.
     Don't instantiate directly, call Var.aux() instead.
nnf.all\_models(names: Iterable[Hashable]) \rightarrow Iterator[Dict[Hashable, bool]]
     Yield dictionaries with all possible boolean values for the names.
     >>> list(all_models(["a", "b"]))
      [{'a': False, 'b': False}, {'a': False, 'b': True}, ...
nnf. complete_models (models: Iterable[Dict[Hashable, bool]], names: Iterable[Hashable]) \rightarrow Itera-
                           tor[Dict[Hashable, bool]]
nnf.decision(var:
                              nnf.Var,
                                           if_true:
                                                          T_NNF,
                                                                      if_false:
                                                                                       U NNF)
                 nnf.Or[typing.Union[nnf.And[typing.Union[nnf.Var,
                                                                                            ~T NNF]],
                                                      ~U_NNF]]]][Union[nnf.And[typing.Union[nnf.Var,
                 nnf.And[typing.Union[nnf.Var,
                 ~T_NNF]][Union[nnf.Var,
                                                                         nnf.And[typing.Union[nnf.Var,
                                                      T_NNF]],
                  ~U NNF]][Union[nnf.Var, U NNF]]]]
     Create a decision node with a variable and two branches.
          Parameters
                 • var – The variable node to decide on.
                 • if true – The branch if the decision is true.
                 • if false – The branch if the decision is false.
```

#### nnf.true = true

A node that's always true. Technically an And node without children.

#### nnf.false = false

A node that's always false. Technically an Or node without children.

#### 2.2 Submodules

## 2.3 nnf.operators module

Convenience functions for logical relationships that are not part of NNF.

These functions will simulate those relationships, often by doubling sentences or altering their structure to negate them. This makes them inefficient.

- $nnf.operators.xor(a:nnf.NNF,b:nnf.NNF) \rightarrow nnf.Or[nnf.And[nnf.NNF]][nnf.And[nnf.NNF]][nnf.NNF]]$ Exactly one of the operands is true.
- nnf.operators.nand (a: nnf.NNF, b: nnf.NNF)  $\rightarrow$  nnf.Or[nnf.NNF][nnf.NNF] At least one of the operands is false.
- $nnf.operators.nor(a: nnf.NNF, b: nnf.NNF) \rightarrow nnf.And[nnf.NNF][nnf.NNF]$ Both of the operands are false.
- nnf.operators.implies (a: nnf.NNF, b: nnf.NNF)  $\rightarrow nnf.Or[nnf.NNF][nnf.NNF]$  b is true whenever a is true.
- nnf.operators.implied\_by (a: nnf.NNF, b: nnf.NNF)  $\rightarrow$  nnf.Or[nnf.NNF][nnf.NNF] a is true whenever b is true.
- $nnf.operators.iff(a: nnf.NNF, b: nnf.NNF) \rightarrow nnf.Or[nnf.And[nnf.NNF]][nnf.And[nnf.NNF]][nnf.NNF]]$  a is true if and only if b is true.
- nnf.operators.and\_(a:  $T\_NNF$ , b:  $U\_NNF$ )  $\rightarrow$  nnf.And[typing.Union[~T\_NNF, ~U\_NNF]] a and b are both true. Included for completeness.
- nnf.operators.or\_(a:  $T\_NNF$ , b:  $U\_NNF$ )  $\rightarrow$  nnf.Or[typing.Union[~T\_NNF, ~U\_NNF]][Union[T\_NNF, U\_NNF]] a or b is true. Included for completeness.

### 2.4 nnf.dimacs module

A parser and serializer for the DIMACS CNF and SAT formats.

nnf.dimacs.dump (obj: nnf.NNF, fp: TextIO, \*, num\_variables: Optional[int] = None, var\_labels: Optional[Dict[Hashable, int]] = None, comment\_header: Optional[str] = None, mode:  $str = 'sat') \rightarrow None$ Dump a sentence into an open file in a DIMACS format.

Variable names have to be integers. If the variables in the sentence you want to dump are not integers, you can pass a var labels dictionary to map names to integers.

#### **Parameters**

- **obj** The sentence to dump.
- **fp** The open file.
- num\_variables Override the number of variables, in case there are variables that don't appear in the sentence.
- var\_labels A dictionary mapping variable names to integers, to rename non-integer variables.
- **comment\_header** A comment to include at the top of the file. May include newlines.
- mode Either 'sat' to dump in the general SAT format, or 'cnf' to dump in the specialized CNF format.
- nnf.dimacs.load (fp: TextIO)  $\rightarrow$  Union[nnf.NNF, nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]]] Load a sentence from an open file.

The format is automatically detected.

2.4. nnf.dimacs module

```
nnf.dimacs.dumps(obj: nnf.NNF, *, num_variables: Optional[int] = None, var_labels: Optional[Dict[Hashable, int]] = None, comment_header: Optional[str] = None, mode: str = 'sat') → str

Like dump(), but to a string instead of to a file.

nnf.dimacs.loads(s: str) → Union[nnf.NNF, nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]]]

Like load(), but from a string instead of from a file.

exception nnf.dimacs.DimacsError

Bases: Exception

exception nnf.dimacs.DimacsError

Bases: nnf.dimacs.DimacsError

exception nnf.dimacs.DecodeError

Bases: nnf.dimacs.DimacsError
```

### 2.5 nnf.dsharp module

Interoperability with DSHARP.

load and loads can be used to parse files created by DSHARP's -Fnnf option.

compile invokes DSHARP directly to compile a sentence. This requires having DSHARP installed.

The parser was derived by studying DSHARP's output and source code. This format might be some sort of established standard, in which case this parser might reject or misinterpret some valid files in the format.

DSHARP may not work properly for some (usually trivially) unsatisfiable sentences, incorrectly reporting there's a solution. This bug dates back to sharpSAT, on which DSHARP was based:

https://github.com/marcthurley/sharpSAT/issues/5

It was independently discovered by hypothesis during testing of this module.

nnf.dsharp.load (fp: TextIO, var\_labels: Optional[Dict[int, Hashable]] = None)  $\rightarrow$  nnf.NNF Load a sentence from an open file.

An optional var\_labels dictionary can map integers to other names.

nnf.dsharp.loads (s: str, var\_labels: Optional[Dict[int, Hashable]] = None)  $\rightarrow$  nnf.NNF Load a sentence from a string.

nnf.dsharp.compile (sentence: nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]][nnf.Var]], executable: str = 'dsharp',  $smooth: bool = False, timeout: Optional[int] = None, extra_args: Sequence[str] = ()) <math>\rightarrow$  nnf.NNF

Run DSHARP to compile a CNF sentence to (s)d-DNNF.

This requires having DSHARP installed.

The returned sentence will be marked as deterministic.

#### **Parameters**

- **sentence** The CNF sentence to compile.
- **executable** The path of the dsharp executable. If the executable is in your PATH there's no need to set this.
- **smooth** Whether to produce a smooth sentence.
- timeout Tell DSHARP to give up after a number of seconds.
- extra\_args Extra arguments to pass to DSHARP.

### 2.6 nnf.amc module

An implementation of algebraic model counting.

```
nnf.amc.eval (node: nnf.NNF, add: Callable[[T, T], T], mul: Callable[[T, T], T], add_neut: T, mul_neut: T, labeling: Callable[[nnf.Var], T]) \to T Execute an AMC technique, given a semiring and a labeling function.
```

#### **Parameters**

- node The sentence to calculate the value of.
- add The operator, to combine nnf. Or nodes.
- mul The operator, to combine nnf.And nodes.
- add  $neut e^{\wedge}$ , the neutral element of the operator.
- $mul_neut e^{\wedge}$ , the neutral element of the operator.
- labeling The labeling function, to assign a value to each variable node.

```
nnf.amc.reduce (node: nnf.NNF, add_key: Optional[Callable[[T], Any]], mul: Callable[[T, T], T], add_neut: T, mul_neut: T, labeling: Callable[[nnf.Var], T]) \rightarrow nnf.NNF Execute AMC reduction on a sentence.
```

In AMC reduction, the operator must be max on some total order, and the branches of the sentence that don't contribute to the maximum value are removed. This leaves a simpler sentence with only the models with a maximum value.

#### **Parameters**

- node The sentence.
- add\_key A function given to max's key argument to determine the total order of the operator. Pass None to use the default ordering.
- **mul** See *eval* ().
- add neut See eval().
- mul\_neut See eval().
- labeling See eval().

**Returns** The transformed sentence.

```
nnf.amc.SAT(node: nnf.NNF) \rightarrow bool
```

Determine whether a DNNF sentence is satisfiable.

```
nnf.amc.NUM\_SAT (node: nnf.NNF) \rightarrow int
```

Determine the number of models that satisfy a sd-DNNF sentence.

nnf .amc . WMC (node: nnf.NNF, weights: Callable[[nnf.Var], float])  $\rightarrow$  float Model counting of sd-DNNF sentences, weighted by variables.

#### **Parameters**

- **node** The sentence to measure.
- weights A dictionary mapping variable nodes to weights.
- nnf.amc.PROB (node: nnf.NNF, probs: Dict[Hashable, float])  $\rightarrow$  float Model counting of d-DNNF sentences, weighted by probabilities.

#### **Parameters**

2.6. nnf.amc module

- node The sentence to measure.
- probs A dictionary mapping variable names to probabilities.

```
nnf.amc.GRAD (node: nnf.NNF, probs: Dict[Hashable, float], k: Optional[Hashable] = None) \rightarrow Tuple[float, float]
```

Calculate a gradient of a d-DNNF sentence being true depending on the value of a variable, given probabilities for all variables.

#### **Parameters**

- **node** The sentence.
- **probs** A dictionary mapping variable names to probabilities.
- **k** The name of the variable to check relative to.

**Returns** A tuple of two floats (probability, gradient).

```
nnf.amc.MPE (node: nnf.NNF, probs: Dict[Hashable, float]) \rightarrow float
```

```
nnf.amc.maxplus_reduce (node: nnf.NNF, labels: Dict[nnf.Var, float]) \rightarrow nnf.NNF Execute AMC reduction using the maxplus algebra.
```

#### **Parameters**

- **node** The sentence.
- labels A dictionary mapping variable nodes to numbers.

### 2.7 nnf.tseitin module

Transformations using the well-known Tseitin encoding.

The Tseitin transformation converts any arbitrary circuit to one in CNF in polynomial time/space. It does so at the cost of introducing new variables (one for each logical connective in the formula).

```
\label{eq:conf_conf} \begin{split} &\text{nnf.tseitin.to\_CNF} \ (\textit{theory:} & \textit{nnf.NNF}, & \textit{simplify:} & \textit{bool} & = & \textit{True}) \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ & & \\ & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ &
```

#### **Parameters**

- theory Theory to convert.
- simplify If True, simplify clauses even if that means eliminating variables.

#### 2.8 nnf.kissat module

Interoperability with kissat.

solve invokes the SAT solver directly on the given theory.

```
nnf.kissat.solve (sentence: nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]][nnf.Var]], extra_args: Sequence[str] = ()) <math>\rightarrow Optional[Dict[Hashable, bool]] Run kissat to compute a satisfying assignment.
```

#### **Parameters**

- sentence The CNF sentence to solve.
- extra\_args Extra arguments to pass to kissat.

## 2.9 nnf.pysat module

- nnf.pysat.satisfiable (sentence:  $nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]][nnf.Var]]) <math>\rightarrow$  bool Return whether a CNF sentence is satisfiable.
- $$\label{eq:continuity} \begin{split} &\text{nnf.pysat.solve} \ (\textit{sentence:} & \textit{nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]][nnf.Var]])} \quad \to \quad & \text{Optional[Dict[Hashable, bool]]} \\ &\text{Return a model of a CNF sentence, or None if unsatisfiable.} \end{split}$$
- $\label{eq:continuity} $$ nnf.pysat.models (sentence: nnf.And[nnf.Or[nnf.Var]][nnf.Or[nnf.Var]][nnf.Var]]) $$ \to $$ Iterator[Dict[Hashable, bool]] $$ Yield all models of a CNF sentence.$
- nnf.pysat.available = False

Indicates whether the PySAT library is installed and available for use.

## CHAPTER 3

### Command line interface

Some of python-nnf's functionality is exposed through a command line tool. It can be invoked as pynnf or python3 -m nnf.

## 3.1 SAT solving

pynnf sat tests whether a sentence is satisfiable, while pynnf sharpsat counts how many solutions it has.

Add -v to get extra information about the sentence and the running time.

#### Example:

```
$ pynnf sat uf20-01.cnf
SATISFIABLE
```

Beware that it's much slower than dedicated solvers like MiniSat.

## 3.2 Sentence summary

pynnf info shows basic information about a sentence.

### Examples:

```
$ pynnf info uf20-01.cnf
Sentence is in CNF.
Variables: 20
Size: 360
Clauses: 90
Clause size: 3
$ pynnf info uf100-016.cnf.nnf
Sentence is decomposable.
```

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```
Variables: 97
Size: 109
```

## 3.3 Visualizing sentences

pynnf draw converts sentences to a DOT representation, and either outputs that or feeds it to dot to immediately output an image.

Immediately outputting an image requires having dot installed. It's done when the output file has an image extension, or when a format is passed with the -f flag.

#### Examples:

See pynnf draw --help for more information.

## CHAPTER 4

Caveats

There are a few things to keep in mind when using python-nnf.

## 4.1 Node duplication

If the same node occurs multiple times in a sentence, then it often pays to make sure that it isn't created multiple times.

Here's a (contrived) example of two ways to construct the same sentence:

```
>>> inefficient = And({
...     Or({A, B}),
...     And({A, Or({A, B})}),
... })
>>> dup_node = Or({A, B})
>>> efficient = And({
...     dup_node,
...     And({A, dup_node}),
... })
```

These objects behave identically, but the first one stores the node  $Or(\{A, B\})$  twice, and the other stores it only once. That means the second one uses less memory.

For a lot of sentences, this isn't worth worrying about. But if you have many nodes that occur multiple times, and they descend from nodes that occur multiple times, you may end up using a lot more memory than necessary.

The .object\_count() and .deduplicate() methods exist to diagnose this problem. .object\_count() tells you how many actual objects are used to represent the sentence, and .deduplicate() returns a maximally compact copy.

If .deduplicate() changes the value of .object\_count() a lot then the sentence could benefit from watching out not to create objects multiple times.

```
>>> inefficient.object_count()
6
```

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```
>>> inefficient.deduplicate().object_count()
5
```

In this case the difference is pretty small.

## 4.2 Decomposability and determinism

A lot of methods are much faster to perform on sentences that are decomposable or deterministic, such as model enumeration.

Decomposability is automatically detected.

Determinism is too expensive to automatically detect, but it can give a huge speedup. If you know a sentence to be deterministic, run .mark\_deterministic() to enable the relevant optimizations.

A compiler like DSHARP may be able to convert some sentences into equivalent deterministic decomposable sentences. The output of DSHARP can be loaded using the nnf.dsharp module. Sentences returned by nnf.dsharp.compile() are automatically marked as deterministic.

### 4.3 Other duplication inefficiencies

Even when properly deduplicated, the kind of sentence that's vulnerable to node duplication might still be inefficient to work with for some operations.

A known offender is equality (==). Currently, if two of such sentences are compared that are equal but don't share any objects, it takes a very long time even if both sentences don't have any duplication within themselves.

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## CHAPTER 5

Introduction

Sentences are made up of nodes. To start with, define some variables:

```
>>> from nnf import Var
>>> A, B, C = Var('A'), Var('B'), Var('C')
```

Then, if you want to write the sentence "A or B":

```
>>> from nnf import Or
>>> sentence = Or({A, B})
>>> sentence = A | B # alternative syntax
```

Or "B and not C":

```
>>> from nnf import And
>>> sentence = And({B, ~C})
>>> sentence = B & ~C
```

Of course you can nest these, for more interesting sentences:

```
>>> sentence = Or({And({A, B}), And({~B, C})})
```

You can ask queries, and perform transformations:

```
>>> sentence.decomposable()
True
>>> sentence.smooth()
False
>>> list(sentence.models())
[{'A': True, 'B': True, 'C': True}, {'A': True, 'B': False, ...
>>> new = sentence.condition({'B': True})
>>> new
Or({And({A, true}), And({false, C})})
>>> list(new.models())
[{'A': True, 'C': True}, {'A': True, 'C': False}]
```

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```
>>> new.simplify()
A
```

## CHAPTER 6

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